

Convex Optimization and Applications

4 - Convex Optimization Problems

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Outline

- 1 Definitions
- 2 Local vs. Global optima
- 3 Problem Reformulations
- 4 First-order Optimality Conditions

Definitions

Optimization problem (aka *nonlinear program*)

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && f_0(\mathbf{x}) && \text{(P)} \\ & \text{s.t.} && f_i(\mathbf{x}) \leq 0 && (\forall i \in [m]). \end{aligned}$$

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- The *optimal value* of (P) is

$$p^* = \inf\{f_0(\mathbf{x}) \mid \mathbf{x} \in \mathcal{F}\} \in \mathbb{R} \cup \{-\infty, +\infty\}.$$

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- $p^* = -\infty \iff \exists (\mathbf{x}_i)_{i \in \mathbb{N}}$ with $\mathbf{x}_i \in \mathcal{F}$ and $f_0(\mathbf{x}_i) \rightarrow \infty$; in this case, (P) is *unbounded from below*.

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- When f_0 is constant, we just have to find any $\mathbf{x} \in \mathcal{F}$; In this case (P) is a *feasibility problem*.
- If $\mathbf{x}^* \in \mathcal{F}$ satisfies $f_0(\mathbf{x}^*) = p^*$, we say that \mathbf{x}^* *solves* (P), or that \mathbf{x}^* is a (global) *optimal solution* to (P)

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- $\mathbf{x} \in \mathcal{F}$ is called ϵ -suboptimal if $f_0(\mathbf{x}) \leq p^* + \epsilon$.
- The set of all ϵ -suboptimal solutions is called the ϵ -suboptimal set.

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- The set of all ϵ -suboptimal solutions is called the ϵ -suboptimal set.
- **Remark:** The constraints are understood in the sense of extended functions, i.e., $\mathbf{x} \notin \text{dom } f_i \implies f_i(\mathbf{x}) = \infty$. Hence,

$$\mathcal{F} \subseteq \bigcap_{i=0}^m \text{dom } f_i.$$

Local optimum

Definition (Local optimum).

The vector x is called a *local optimum* for Problem (P) if it solves the problem

$$\begin{aligned} & \underset{z \in \mathbb{R}^n}{\text{minimize}} && f_0(z) && && (\text{P}_R) \\ & \text{s.t.} && f_i(z) \leq 0 && (\forall i \in [m]); \\ & && \|z - x\| \leq R \end{aligned}$$

for some $R > 0$. In other words, there is a neighbourhood of x in \mathcal{F} in which f_0 is minimized at x .

Local optimum

Proposition (Differential characterization of local optima).

Assume the objective function f_0 is twice differentiable, and let $\mathbf{x} \in \text{int } \mathcal{F}$. Then, the following holds:

- If $\nabla f_0(\mathbf{x}) = \mathbf{0}$ and $\nabla^2 f_0(\mathbf{x}) \succ 0$, then \mathbf{x} is a local optimum.
- Conversely, if \mathbf{x} is a local optimum, then $\nabla f_0(\mathbf{x}) = \mathbf{0}$ and $\nabla^2 f_0(\mathbf{x}) \succeq 0$.

Remarks:

- Above proposition is valid **for interior points only** !

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Remarks:

- Above proposition is valid **for interior points only** !
- cannot replace \succ by \succeq in the 1st statement ($f(x) = x^3$)
- cannot replace \succeq by \succ in the 2d statement ($f(x) = x^4$)

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Problem (P) is said to be *convex* if f_0, \dots, f_m are convex.

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Remark: We also say that the maximization problem

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{maximize}} && f_0(\mathbf{x}) \\ & \text{s.t.} && f_i(\mathbf{x}) \geq 0 && (\forall i \in [m]). \end{aligned}$$

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Problem (P) is said to be *convex* if f_0, \dots, f_m are convex.

Equality constraints $f_i(\mathbf{x}) = 0$ can be handled as $f_i(\mathbf{x}) \leq 0, f_i(\mathbf{x}) \geq 0$. Hence, in a convex optimization problem, equality constraints must be *linear*.

It is often convenient to write equality constraints separately:

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && f_0(\mathbf{x}) && \text{(P}_{\text{Eq}}) \\ & \text{s.t.} && f_i(\mathbf{x}) \leq 0 && (\forall i \in [m]); \\ & && A\mathbf{x} = \mathbf{b}. \end{aligned}$$

Fundamental result

Theorem

Let (P) be a convex optimization problem. Then, any local optimum x^* is also a global optimum.

Equivalence of Problems

“informal” definition

We say that Problems (P) and (Q) are equivalent, and we use the (nonstandard) notation $P \sim Q$, if there is a “simple transformation” which maps an optimal solution of P to an optimal solution of Q , and vice-versa.

Change of variables

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && f_0(\mathbf{x}) && \text{(P)} \\ & \text{s.t.} && f_i(\mathbf{x}) \leq 0 && (\forall i \in [m]). \end{aligned}$$

If $\phi : \mathbb{R}^n \rightarrow \mathbb{R}^n$ is one-to-one, and every feasible \mathbf{x} can be written as $\mathbf{x} = \phi(\mathbf{z})$ for some \mathbf{z} , then

$$\begin{aligned} \text{P} & \sim \underset{\mathbf{z} \in \mathbb{R}^n}{\text{minimize}} && f_0(\phi(\mathbf{z})) \\ & \text{s.t.} && f_i(\phi(\mathbf{z})) \leq 0 && (\forall i \in [m]). \end{aligned}$$

Transformation of objective or constraints

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && f_0(\mathbf{x}) && \text{(P)} \\ & \text{s.t.} && f_i(\mathbf{x}) \leq 0 && (\forall i \in [m]). \end{aligned}$$

If

- $\psi_0 : \mathbb{R} \rightarrow \mathbb{R}$ is strictly increasing
- $\forall i \in [m], \psi_i : \mathbb{R} \rightarrow \mathbb{R}$ satisfies $\psi_i(u) \leq 0 \iff u \leq 0$,

then we have:

$$\begin{aligned} \text{P} & \sim \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && \psi_0(f_0(\mathbf{x})) \\ & \text{s.t.} && \psi_i(f_i(\mathbf{x})) \leq 0 && (\forall i \in [m]). \end{aligned}$$

Eliminating equality constraints

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && f_0(\mathbf{x}) && (\mathbf{P}_{\text{Eq}}) \\ & \text{s.t.} && f_i(\mathbf{x}) \leq 0 && (\forall i \in [m]) \\ & && A\mathbf{x} = \mathbf{b}. \end{aligned}$$

The equality constraints define an affine subspace $L = \{\mathbf{x} \in \mathbb{R}^n : A\mathbf{x} = \mathbf{b}\} = \{C\mathbf{z} + \mathbf{d} : \mathbf{z} \in \mathbb{R}^r\}$ for some $C \in \mathbb{R}^{n \times r}$, $\mathbf{d} \in \mathbb{R}^n$. Hence,

$$\mathbf{P}_{\text{Eq}} \quad \sim \quad \underset{\mathbf{z} \in \mathbb{R}^r}{\text{minimize}} \quad f_0(C\mathbf{z} + \mathbf{d}) \\ \text{s.t.} \quad f_i(C\mathbf{z} + \mathbf{d}) \leq 0 \quad (\forall i \in [m]).$$

Slack variables

We can replace linear inequalities by linear equalities by introducing *slack variables*. For example,

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && f_0(\mathbf{x}) \\ & \text{s.t.} && f_i(\mathbf{x}) \leq 0 \quad (\forall i \in [m]) \\ & && A\mathbf{x} \leq \mathbf{b}, \end{aligned}$$

where $A \in \mathbb{R}^{p \times n}$, is equivalent to

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n, \mathbf{s} \in \mathbb{R}^p}{\text{minimize}} && f_0(\mathbf{x}) \\ & \text{s.t.} && f_i(\mathbf{x}) \leq 0 \quad (\forall i \in [m]) \\ & && A\mathbf{x} + \mathbf{s} = \mathbf{b} \\ & && \mathbf{s} \geq \mathbf{0}. \end{aligned}$$

Epigraph formulation

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && f_0(\mathbf{x}) && \text{(P)} \\ & \text{s.t.} && f_i(\mathbf{x}) \leq 0 && (\forall i \in [m]). \end{aligned}$$

We can always assume that the objective function is a linear form: $f_0(\mathbf{x}) = \mathbf{c}^T \mathbf{x}$, with $\|\mathbf{c}\| = 1$, by passing to the *epigraph formulation*:

$$\begin{aligned} \text{P} & \sim \underset{\mathbf{x} \in \mathbb{R}^n, t \in \mathbb{R}}{\text{minimize}} && t \\ & \text{s.t.} && f_i(\mathbf{x}) \leq 0 && (\forall i \in [m]) \\ & && f_0(\mathbf{x}) \leq t. \end{aligned}$$

Partial minimization

It is possible to reduce a problem by solving it (partially) for some blocks of variables. For example, the following two problems are equivalent:

$$\begin{aligned} & \underset{\mathbf{x}_1 \in \mathbb{R}^{n_1}, \mathbf{x}_2 \in \mathbb{R}^{n_2}}{\text{minimize}} && f_0(\mathbf{x}_1, \mathbf{x}_2) \\ & \text{s.t.} && f_i(\mathbf{x}_1) \leq 0 \quad (\forall i \in [m_1]) \\ & && g_j(\mathbf{x}_1, \mathbf{x}_2) \leq 0 \quad (\forall j \in [m_2]) \end{aligned}$$

$$\begin{aligned} \sim & \underset{\mathbf{x}_1 \in \mathbb{R}^{n_1}}{\text{minimize}} && \tilde{f}_0(\mathbf{x}_1) \\ & \text{s.t.} && f_i(\mathbf{x}_1) \leq 0 \quad (\forall i \in [m_1]), \end{aligned}$$

where we have defined

$$\tilde{f}_0(\mathbf{x}_1) := \inf \{ f_0(\mathbf{x}_1, \mathbf{x}_2) \mid \mathbf{x}_2 \in \mathbb{R}^{n_2}, g_j(\mathbf{x}_1, \mathbf{x}_2) \leq 0, \forall j \in [m_2] \}.$$

First-order optimality condition

This result gives a simple optimality condition, which depends only on ∇f_0 and the feasibility set

$$\mathcal{F} = \{\mathbf{x} \in \mathbb{R}^n \mid f_1(\mathbf{x}) \leq 0, \dots, f_m(\mathbf{x}) \leq 0\}.$$

Theorem

Let f_0 be differentiable. Then, a vector \mathbf{x}^* is optimal for the convex problem (P) if and only if \mathbf{x}^* is feasible (i.e., $\mathbf{x}^* \in \mathcal{F}$), and

$$\forall \mathbf{y} \in \mathcal{F}, \quad \nabla f_0(\mathbf{x}^*)^T (\mathbf{y} - \mathbf{x}^*) \geq 0.$$

Geometrically, this means that either $\nabla f_0(\mathbf{x}^*) = \mathbf{0}$, or $\nabla f_0(\mathbf{x}^*)$ defines a supporting hyperplane to \mathcal{F} at \mathbf{x}^* .

Equality-constrained optimization

We can use the previous theorem to characterize optimal solutions of equality-constrained convex programs:

Proposition

Consider the optimization problem

$$\begin{array}{ll} \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} & f_0(\mathbf{x}) \\ \text{s.t.} & A\mathbf{x} = \mathbf{b}, \end{array}$$

where f_0 is convex and differentiable. Then, \mathbf{x}^* is optimal iff $\nabla f_0(\mathbf{x}^*) \in \mathbf{Im} A^T$.

Optimization over \mathbb{R}_+^n

We can also obtain a characterization of optimal solutions of convex programs over the nonnegative orthant:

Proposition

Consider the optimization problem

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && f_0(\mathbf{x}) \\ & \text{s.t.} && \mathbf{x} \geq \mathbf{0}, \end{aligned}$$

where f_0 is convex and differentiable. Then, \mathbf{x}^* is optimal iff

- $\mathbf{x}^* \geq \mathbf{0}$
- $\nabla f_0(\mathbf{x}^*) \geq \mathbf{0}$
- $\forall i \in [n], \left(x_i = 0 \text{ or } \frac{\partial f_0}{\partial x_i}(\mathbf{x}^*) = 0 \right)$