

# Convex Optimization and Applications

## 4 - Convex Optimization Problems

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# Outline

- 1 Definitions
- 2 Local vs. Global optima
- 3 Problem Reformulations
- 4 First-order Optimality Conditions

# Definitions

Optimization problem (aka *nonlinear program*)

$$\begin{array}{ll} \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} & f_0(\mathbf{x}) \\ \text{s.t.} & f_i(\mathbf{x}) \leq 0 \quad (\forall i \in [m]). \end{array} \quad (\text{P})$$

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- The *optimal value* of (P) is

$$p^* = \inf\{f_0(\mathbf{x}) \mid \mathbf{x} \in \mathcal{F}\} \in \mathbb{R} \cup \{-\infty, +\infty\}.$$

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- When  $f_0$  is constant, we just have to find any  $x \in \mathcal{F}$ ; In this case (P) is a *feasibility problem*.
- If  $x^* \in \mathcal{F}$  satisfies  $f_0(x^*) = p^*$ , we say that  $x^*$  *solves* (P), or that  $x^*$  is a (global) *optimal solution* to (P)

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## Definition

- $\mathbf{x} \in \mathcal{F}$  is called  $\epsilon$ -suboptimal if  $f_0(\mathbf{x}) \leq p^* + \epsilon$ .
- The set of all  $\epsilon$ -suboptimal solutions is called the  $\epsilon$ -suboptimal set.

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- The set of all  $\epsilon$ -suboptimal solutions is called the  $\epsilon$ -suboptimal set.
- **Remark:** The constraints are understood in the sense of extended functions, i.e.,  $\mathbf{x} \notin \text{dom } f_i \implies f_i(\mathbf{x}) = \infty$ . Hence,

$$\mathcal{F} \subseteq \bigcap_{i=0}^m \text{dom } f_i.$$

# Local optimum

## Definition (Local optimum).

The vector  $x$  is called a *local optimum* for Problem (P) if it solves the problem

$$\begin{aligned} & \underset{z \in \mathbb{R}^n}{\text{minimize}} && f_0(z) && && (\text{P}_R) \\ & \text{s.t.} && f_i(z) \leq 0 && (\forall i \in [m]); \\ & && \|z - x\| \leq R \end{aligned}$$

for some  $R > 0$ . In other words, there is a neighbourhood of  $x$  in  $\mathcal{F}$  in which  $f_0$  is minimized at  $x$ .

# Local optimum

## Proposition (Differential characterization of local optima).

Assume the objective function  $f_0$  is twice differentiable, and let  $x \in \text{int } \mathcal{F}$ . Then, the following holds:

- If  $\nabla f_0(x) = \mathbf{0}$  and  $\nabla^2 f_0(x) \succ 0$ , then  $x$  is a local optimum.
- Conversely, if  $x$  is a local optimum, then  $\nabla f_0(x) = \mathbf{0}$  and  $\nabla^2 f_0(x) \succeq 0$ .

### Remarks:

- Above proposition is valid **for interior points only** !

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### Remarks:

- Above proposition is valid **for interior points only** !
- can't replace  $\succ$  by  $\succeq$  in the 1st statement ( $f(x) = x^3$ )
- can't replace  $\succeq$  by  $\succ$  in the 2d statement ( $f(x) = x^4$ )



# Convex optimization problem

Optimization problem

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Problem (P) is said to be *convex* if  $f_0, \dots, f_m$  are convex.

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Problem (P) is said to be *convex* if  $f_0, \dots, f_m$  are convex.

**Remark:** We also say that the maximization problem

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{maximize}} && f_0(\mathbf{x}) \\ & \text{s.t.} && f_i(\mathbf{x}) \geq 0 && (\forall i \in [m]). \end{aligned}$$

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Problem (P) is said to be *convex* if  $f_0, \dots, f_m$  are convex.

Equality constraints  $f_i(\mathbf{x}) = 0$  can be handled as  $f_i(\mathbf{x}) \leq 0, f_i(\mathbf{x}) \geq 0$ . Hence, in a convex optimization problem, equality constraints must be *linear*.

It is often convenient to write equality constraints separately:

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && f_0(\mathbf{x}) && \text{(P}_{\text{Eq}}) \\ & \text{s.t.} && f_i(\mathbf{x}) \leq 0 && (\forall i \in [m]); \\ & && A\mathbf{x} = \mathbf{b}. \end{aligned}$$

# Fundamental result

## Theorem

Let (P) be a convex optimization problem. Then, any local optimum  $x^*$  is also a global optimum.

# Equivalence of Problems

## “informal” definition

We say that Problems (P) and (Q) are equivalent, and we use the (nonstandard) notation  $P \sim Q$ , if there is a “simple transformation” which maps an optimal solution of  $P$  to an optimal solution of  $Q$ , and vice-versa.

# Change of variables

$$\begin{aligned} & \underset{x \in \mathbb{R}^n}{\text{minimize}} && f_0(x) && \text{(P)} \\ & \text{s.t.} && f_i(x) \leq 0 && (\forall i \in [m]). \end{aligned}$$

If  $\phi : \mathbb{R}^n \rightarrow \mathbb{R}^n$  is one-to-one, and every feasible  $x$  can be written as  $x = \phi(z)$  for some  $z$ , then

$$\begin{aligned} \text{P} & \sim \underset{z \in \mathbb{R}^n}{\text{minimize}} && f_0(\phi(z)) \\ & \text{s.t.} && f_i(\phi(z)) \leq 0 && (\forall i \in [m]). \end{aligned}$$

# Transformation of objective or constraints

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && f_0(\mathbf{x}) && \text{(P)} \\ & \text{s.t.} && f_i(\mathbf{x}) \leq 0 && (\forall i \in [m]). \end{aligned}$$

If

- $\psi_0 : \mathbb{R} \rightarrow \mathbb{R}$  is strictly increasing
- $\forall i \in [m], \psi_i : \mathbb{R} \rightarrow \mathbb{R}$  satisfies  $\psi_i(u) \leq 0 \iff u \leq 0$ ,

then we have:

$$\begin{aligned} \text{P} & \sim \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && \psi_0(f_0(\mathbf{x})) \\ & \text{s.t.} && \psi_i(f_i(\mathbf{x})) \leq 0 && (\forall i \in [m]). \end{aligned}$$

# Eliminating equality constraints

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && f_0(\mathbf{x}) && (\mathbf{P}_{\text{Eq}}) \\ & \text{s.t.} && f_i(\mathbf{x}) \leq 0 && (\forall i \in [m]) \\ & && A\mathbf{x} = \mathbf{b}. \end{aligned}$$

The equality constraints define an affine subspace  $L = \{\mathbf{x} \in \mathbb{R}^n : A\mathbf{x} = \mathbf{b}\} = \{C\mathbf{z} + \mathbf{d} : \mathbf{z} \in \mathbb{R}^r\}$  for some  $C \in \mathbb{R}^{n \times r}$ ,  $\mathbf{d} \in \mathbb{R}^n$ . Hence,

$$\mathbf{P}_{\text{Eq}} \quad \sim \quad \underset{\mathbf{z} \in \mathbb{R}^r}{\text{minimize}} \quad f_0(C\mathbf{z} + \mathbf{d}) \\ \text{s.t.} \quad f_i(C\mathbf{z} + \mathbf{d}) \leq 0 \quad (\forall i \in [m]).$$



# Slack variables

We can replace linear inequalities by linear equalities by introducing *slack variables*. For example,

$$\begin{aligned} & \underset{x \in \mathbb{R}^n}{\text{minimize}} && f_0(x) \\ & \text{s.t.} && f_i(x) \leq 0 \quad (\forall i \in [m]) \\ & && Ax \leq b, \end{aligned}$$

where  $A \in \mathbb{R}^{p \times n}$ , is equivalent to

$$\begin{aligned} & \underset{x \in \mathbb{R}^n, s \in \mathbb{R}^p}{\text{minimize}} && f_0(x) \\ & \text{s.t.} && f_i(x) \leq 0 \quad (\forall i \in [m]) \\ & && Ax + s = b \\ & && s \geq \mathbf{0}. \end{aligned}$$

# Epigraph formulation

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && f_0(\mathbf{x}) && \text{(P)} \\ & \text{s.t.} && f_i(\mathbf{x}) \leq 0 && (\forall i \in [m]). \end{aligned}$$

We can always assume that the objective function is a linear form:  $f_0(\mathbf{x}) = \mathbf{c}^T \mathbf{x}$ , with  $\|\mathbf{c}\| = 1$ , by passing to the *epigraph formulation*:

$$\begin{aligned} \text{P} & \sim \underset{\mathbf{x} \in \mathbb{R}^n, t \in \mathbb{R}}{\text{minimize}} && t \\ & \text{s.t.} && f_i(\mathbf{x}) \leq 0 && (\forall i \in [m]) \\ & && f_0(\mathbf{x}) \leq t. \end{aligned}$$

# Partial minimization

It is possible to reduce a problem by solving it (partially) for some blocks of variables. For example, the following two problems are equivalent:

$$\begin{aligned} & \underset{\mathbf{x}_1 \in \mathbb{R}^{n_1}, \mathbf{x}_2 \in \mathbb{R}^{n_2}}{\text{minimize}} && f_0(\mathbf{x}_1, \mathbf{x}_2) \\ & \text{s.t.} && f_i(\mathbf{x}_1) \leq 0 \quad (\forall i \in [m_1]) \\ & && g_j(\mathbf{x}_1, \mathbf{x}_2) \leq 0 \quad (\forall j \in [m_2]) \end{aligned}$$

$$\begin{aligned} \sim & \underset{\mathbf{x}_1 \in \mathbb{R}^{n_1}}{\text{minimize}} && \tilde{f}_0(\mathbf{x}_1) \\ & \text{s.t.} && f_i(\mathbf{x}_1) \leq 0 \quad (\forall i \in [m_1]), \end{aligned}$$

where we have defined

$$\tilde{f}_0(\mathbf{x}_1) := \inf \left\{ f_0(\mathbf{x}_1, \mathbf{x}_2) \mid \mathbf{x}_2 \in \mathbb{R}^{n_2}, g_j(\mathbf{x}_1, \mathbf{x}_2) \leq 0, \forall j \in [m_2] \right\}.$$

# First-order optimality condition

This result gives a simple optimality condition, which depends only on  $\nabla f_0$  and the feasibility set

$$\mathcal{F} = \{\mathbf{x} \in \mathbb{R}^n \mid f_1(\mathbf{x}) \leq 0, \dots, f_m(\mathbf{x}) \leq 0\}.$$

## Theorem

Let  $f_0$  be differentiable. Then, a vector  $\mathbf{x}^*$  is optimal for the convex problem (P) if and only if  $\mathbf{x}^*$  is feasible (i.e.,  $\mathbf{x}^* \in \mathcal{F}$ ), and

$$\forall \mathbf{y} \in \mathcal{F}, \quad \nabla f_0(\mathbf{x}^*)^T (\mathbf{y} - \mathbf{x}^*) \geq 0.$$

Geometrically, this means that either  $\nabla f_0(\mathbf{x}^*) = \mathbf{0}$ , or  $\nabla f_0(\mathbf{x}^*)$  defines a supporting hyperplane to  $\mathcal{F}$  at  $\mathbf{x}^*$ .

# Equality-constrained optimization

We can use the previous theorem to characterize optimal solutions of equality-constrained convex programs:

## Proposition

Consider the optimization problem

$$\begin{array}{ll} \underset{x \in \mathbb{R}^n}{\text{minimize}} & f_0(x) \\ \text{s.t.} & Ax = b, \end{array}$$

where  $f_0$  is convex and differentiable. Then,  $x^*$  is optimal iff  $\nabla f_0(x^*) \in \mathbf{Im} A^T$ .

# Optimization over $\mathbb{R}_+^n$

We can also obtain a characterization of optimal solutions of convex programs over the nonnegative orthant:

## Proposition

Consider the optimization problem

$$\begin{aligned} & \underset{x \in \mathbb{R}^n}{\text{minimize}} && f_0(x) \\ & \text{s.t.} && x \geq \mathbf{0}, \end{aligned}$$

where  $f_0$  is convex and differentiable. Then,  $x^*$  is optimal iff

- $x^* \geq \mathbf{0}$
- $\nabla f_0(x^*) \geq \mathbf{0}$
- $\forall i \in [n], \left( x_i = 0 \text{ or } \frac{\partial f_0}{\partial x_i}(x^*) = 0 \right)$