

6. Exercise sheet

Due-date: Friday, 30.05.2014, *before* the lecture has started

This is the last exercise sheet in the first half of the semester.

Exercise 26

4 points

Prove Birkhoff's Theorem by transforming the minimum-weight perfect matching problem in a bipartite graph to a minimum-cost flow problem.

Exercise 27

4 points

Let $n \in \mathbb{Z}_{>0}$. A matrix $A \in [0, 1]^{n \times n}$ is *doubly stochastic* if the entries in each row and in each column sum up to 1. Show that each doubly stochastic matrix is a convex combination of permutation matrices.

Exercise 28

4 points

Let $G = (V, E)$ be an undirected graph and $c \in \mathbb{R}^E$. Consider the linear programming relaxation of the minimum-weight perfect matching problem:

$$\begin{aligned} \min \quad & \sum_{e \in E} c_e \cdot x_e \\ \text{s.t.} \quad & x(\delta(v)) = 1 \quad \text{for all } v \in V \\ & x \geq 0 \end{aligned}$$

Show that there exists a feasible LP solution if and only if there exists a subset of edges and a set of odd circuits in G such that every node of G is a node of exactly one of the odd circuits or is incident to exactly one of the edges, but not both.

Exercise 29

4 points

Prove that the optimal value of the following problem is the minimum weight of a perfect matching of G .

$$\begin{aligned} \min \quad & \sum_{e \in E} c_e \cdot x_e \\ \text{s.t.} \quad & x(\delta(v)) = 1 \quad \text{for all } v \in V \\ & x(\gamma(S)) \leq (|S| - 1)/2 \quad \text{for all } S \subseteq V, |S| \geq 3, |S| \text{ odd} \\ & x \geq 0 \end{aligned}$$

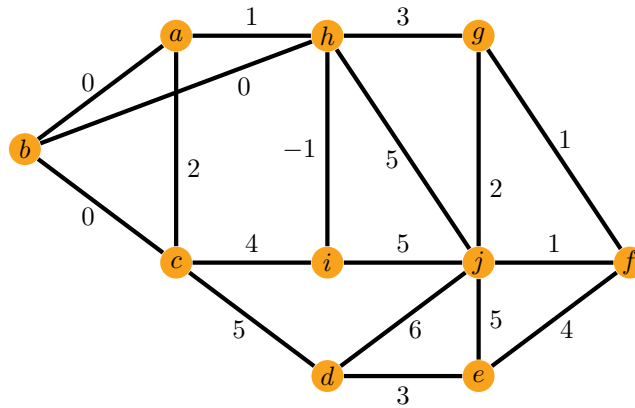
Exercise 30

4 points

Let M be a perfect matching of G and define the "cost" of an M -alternating circuit C of G to be $c(E(C) \setminus M) - c(E(C) \cap M)$. Prove that M is of minimum weight with respect to c if and only if there is no M -alternating circuit of negative cost.

Exercise 31**Tutorial session – 0 points**

Find a minimum-weight perfect matching and an optimal solution of the dual problem for the following weighted graph:

**Exercise 32****Tutorial session – 0 points**

Let $G = (V, E)$ be an undirected graph, let $c \in \mathbb{R}^E$, and let k be a positive integer. Show how the problem of finding a minimum-weight matching having cardinality k can be reduced to a minimum-weight perfect matching problem.

Exercise 33**Tutorial session – 0 points**

Show that there exists a non-bipartite graph with the property that, for every $c \in \mathbb{R}^E$, the optimal value of the linear program in Exercise 32 is equal to the minimum weight of a perfect matching of G .

Exercise 34**Tutorial session – 0 points**

Let $G = (V, E)$ be an undirected graph. Prove that G has a perfect matching if and only if there is an $x \in \mathbb{R}^E$ with:

$$\begin{aligned}
 x(\delta(v)) &= 1 && \text{for all } v \in V \\
 x(\delta(S)) &\geq 1 && \text{for all } S \subseteq V, |S| \text{ odd} \\
 x &\geq 0
 \end{aligned}$$