

1. Exercise Session

23.04.14.

• Orga, Fourier-Motzkin, Farkas

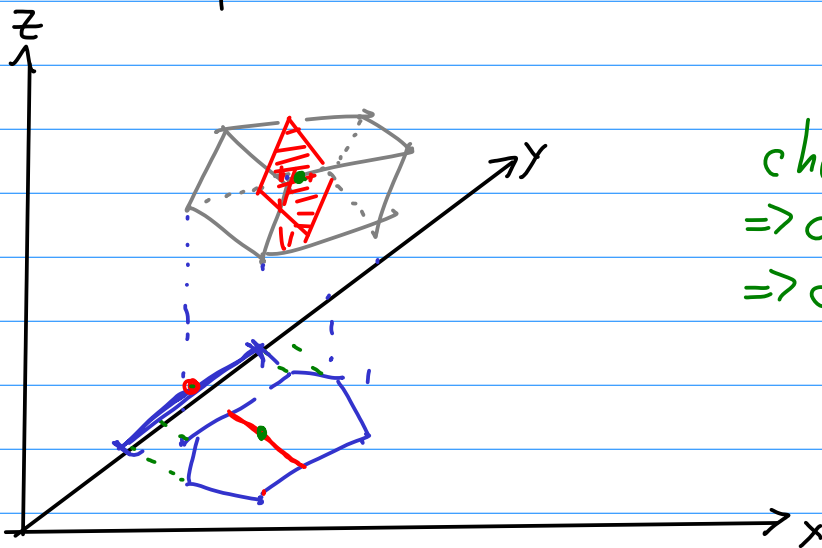
Torsten

geller@math.tu-berlin.de
314 78796

#AO: noch keine Gruppe \Rightarrow Email mit Nr. 0

Fourier Motzkin Elimination:

- procedure to solve question $\exists x \in P$ for $P := \{x \in \mathbb{R}^n \mid Ax \leq b\}$
 - very inefficient: expon. running time
 - nice theoretical tool: derivation of the Farkas Lemma
- idea: project the polytope onto a smaller dimension to get rid of a variable



choose $y \in$ Interval
 \Rightarrow choice for x
 \Rightarrow choice for z

formula: $Ax \leq b$ $x' := (x_2, \dots, x_n), A' := (A_2 \mid \dots \mid A_n)$

$$Ax \leq b \Leftrightarrow \begin{aligned} x_1 + a'_i x' &\leq b_i; & \forall i \in M^+ \\ -x_1 + a'_i x' &\leq b_i; & \forall i \in M^- \\ a'_i x' &\leq b_i; & \forall i \in M \end{aligned}$$

$$\Leftrightarrow \max_{i \in M^-} (a'_i x' - b_i) \leq x_1 \leq \min_{i \in M^+} (b_i - a'_i x')$$
$$a'_i x' \leq b_i \quad \forall i \in M$$

$$\Leftrightarrow \begin{aligned} & \cdot a_i^+ x^1 - b_i \leq b_j - a_j^+ x^1 & \forall i \in M^-, \forall j \in M^+ \\ & a_i^+ x^1 \leq b_i & \forall i \in M \end{aligned}$$

$$\Leftrightarrow \begin{aligned} & (a_i^+ + a_j^+) x^1 \leq (b_j - b_i) & \forall i \in M^-, \forall j \in M^+ \\ & a_i^+ x^1 \leq b_i & \forall i \in M \end{aligned}$$

$$\underbrace{\hspace{10em}}_{=: A'' x^1 \leq b^1}$$

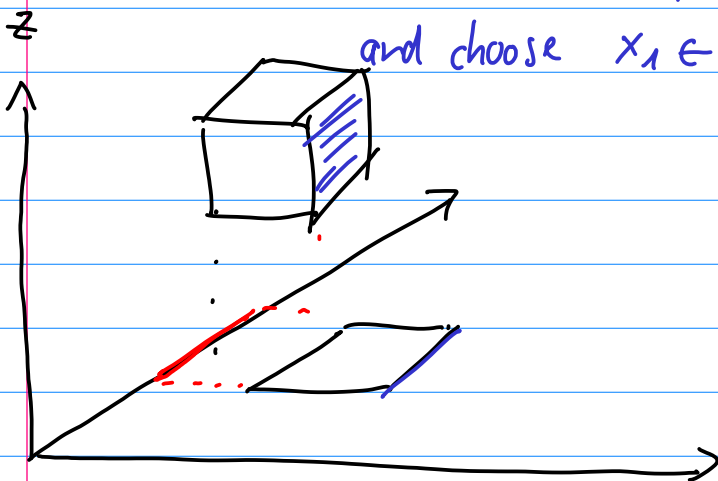
$\Rightarrow \exists \text{ME}$: \cdot find $x^1 = x_2, \dots, x_n$ with $A'' x^1 \leq b^1$
 \cdot choose any value for x_1 from interval $[\max_{i \in M^-} (a_i^+ x^1 - b_i); \min_{i \in M^+} (b_i - a_i^+ x^1)]$

\cdot solving LP: $\min c^T x$
 $Ax \leq b$
 $(x \geq 0)$ \Rightarrow find (x, λ) st. $\tilde{A} x \leq \tilde{b}$
 $c^T x - \lambda \leq 0$
 \cdot eliminate x_1, \dots, x_n
 \cdot choose λ as small as possible
 $\leadsto c^T x = \lambda$

example: cube $1 \leq x_i \leq 2 \quad i \in [n] \quad M: x_i \leq 2 \quad \forall i \in \{2, \dots, n\}$
 $x_i \leq 2 \quad \forall i \in [n] \leadsto -x_i \leq -1$
 $-x_i \leq -1 \quad \forall i \in [n] \quad M^+: x_1 \leq 2$
 $M^-: -x_1 \leq -1$
 $0 \leq 1$

\leadsto find $x_2, \dots, x_n: x_i \leq 2$
 $-x_i \leq -1 \quad \forall i \in \{3, \dots, n\}$

and choose $x_1 \in [1, 2]$



Farkas Lemma:

Farkas Lemma

$A \in \mathbb{R}^{m \times n}, b \in \mathbb{R}^n$. Exactly one of the following alternatives hold:

- 1 $\exists x \geq 0$ such that $Ax = b$ · $p \neq \emptyset$
- 2 \exists some vector p such that $p^T A \geq 0$ and $p^T b < 0$ · $p = \emptyset$

$$P := \{x \mid Ax = b\}$$

$\leadsto p, x$ are certificate for answer if $p = \emptyset$?

$$\Rightarrow LP \in NP, LP \in coNP$$

(Farkas Lemma: Condition for existence of sol.)

Overview:

$$\text{Simplex} \Rightarrow \text{strong duality thm} \Rightarrow \text{Farkas} \Leftarrow \text{FME}$$

Alternative Form:

Alternative Farkas Lemma

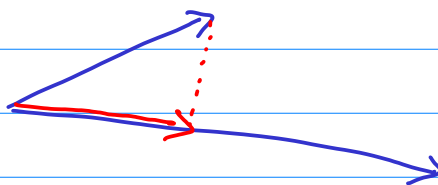
Let $a_1, \dots, a_m \in \mathbb{R}^n, b \in \mathbb{R}^n$. Following statements are equivalent.

- 1 for all $y \in \mathbb{R}^n: y^T a_i \geq 0 \forall i = 1, \dots, m \Rightarrow y^T b \geq 0$

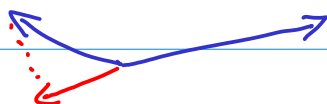
for all y : $\left(\begin{array}{l} y \text{ has a non-neg. projection} \\ \text{onto every } a_i \end{array} \Rightarrow y \text{ has non-neg proj.} \right.$
 $\left. \text{onto } b \right)$

- 2 $b \in \text{Cone}(a_1, \dots, a_m)$ b lies in the cone spanned by a_1, \dots, a_m

non neg. projection:



neg. proj.:

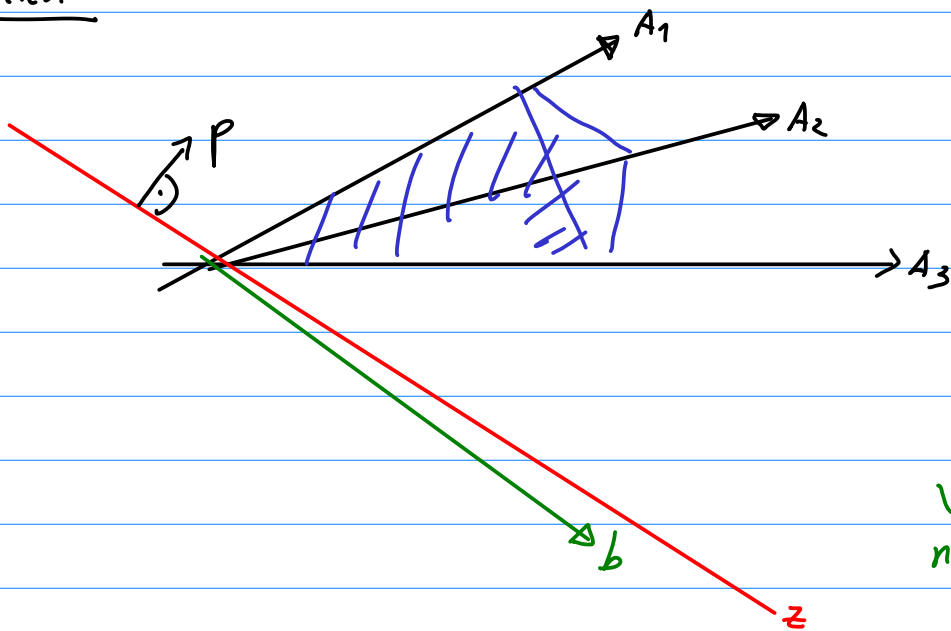


Reformulation: $\neg(1): \exists y \in \mathbb{R}^n: y^T A \geq 0$ and $y^T b < 0$

$\neg(2): Ax = b, x \geq 0$ is not solvable

(can not combine b as a pos. linear comb. of A_i)

Sketch:



$-b \notin \text{Cone}(A_1, A_2, A_3)$
• hyperplane $\{z \mid p^T z = 0\}$ separates b from A_i

• $p^T A \geq 0$
(non-neg. projector)

• $p^T b < 0$ (neg. proj.)

not possible if $b \in \text{Cone}$

Farkas \Rightarrow Strong Duality Theorem

strong. duality thm:

$$(P) \quad \min c^T x \\ Ax = b \\ x \geq 0$$

$$(D) \quad \max y^T b \\ y^T A \leq c^T \\ y \text{ free}$$

Thm: (P) has an opt. sol. if and only if (D) has an optimal solution and their obj. values coincide.

Proof: (via Farkas)

Theorem 9.1.

The system $A \cdot x \leq b$ has a solution x , if and only if there is no vector y satisfying $y \geq 0$, $y^T \cdot A = 0$ and $y^T \cdot b < 0$.

• y^* is opt. solution of (D), $\delta := y^{*T}b$

• $y^*A \leq c^T$
 $-y^*b \leq -\delta - \varepsilon$ has no solution for $\varepsilon > 0$

Farkas

$\Rightarrow \exists \lambda \geq 0: \lambda^T \cdot \begin{pmatrix} -b \\ A \end{pmatrix} = 0$ and $\lambda^T \begin{pmatrix} -\delta - \varepsilon \\ c \end{pmatrix} < 0$

• $\lambda = \begin{pmatrix} \lambda_1 \\ \lambda' \end{pmatrix}$ with $\lambda' \in \mathbb{R}^n$

assume: $\lambda_1 = 0: \lambda' \cdot A = 0$ and $\lambda' \cdot c < 0$,

remember: $A^T y^* \leq c^T \rightarrow$ ~~Farkas~~ Farkas $\Rightarrow \lambda_1 \neq 0$

$\Rightarrow \lambda_1 > 0$, scale $\begin{pmatrix} \lambda_1 \\ \lambda' \end{pmatrix}$ such that $\lambda_1 = 1$

$\lambda^{*T} \cdot A = b^T$ and $\lambda^{*T} c < \delta + \varepsilon$

$A \cdot \lambda' = b$,
feas. sol. for P

$c^T \lambda' < \delta + \varepsilon \Rightarrow$ opt. value for (P) is δ
(via weak duality)

$\Rightarrow \square$

