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## **Compromise Solutions**

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## Motivation

- find *balanced* solutions
- reference point methods: part of many state-of-the-art MCDM tools, little theoretical background











#### Definition (Ideal Point)

Given a multicriteria optimization problem  $\max_{y \in \mathcal{Y}} y$ , the *ideal point*  $y^* = (y_1^*, \dots, y_k^*)$  is defined by

$$y_i^* = \max_{y \in \mathcal{Y}} y_i \qquad \forall \ i.$$



#### Definition (Compromise Solution, Yu 1973)

Given a multicriteria optimization problem  $\max_{y \in \mathcal{Y}} y$  with the ideal point  $y^* \in \mathbb{Q}^k$ , the compromise colution w.r.t. the norm  $\|\cdot\|$  on  $\mathbb{Q}^k$  is

$$y^{\mathsf{cs}} = \min_{y \in \mathcal{Y}} \|y^* - y\|.$$



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 $p = \infty$ 

#### The norms we consider:

 $\ell^p$ -Norm

$$\|y\|_{p} \coloneqq \left(\sum_{i=1}^{k} y_{i}^{p}\right)^{1/p}, \quad p \in [1, \infty) \qquad (\ell^{p}\text{-Norm})$$
$$\|y\|_{\infty} \coloneqq \max_{i=1,\dots,k} y_{i} \qquad (Maximum \ (\ell^{\infty}\text{-})\text{Norm})$$
$$\|y\|_{p} \coloneqq \|y\|_{\infty} + \frac{1}{p} \|y\|_{1}, \quad p \in [1, \infty] \qquad (Cornered \ p\text{-Norm})$$

Degree of balancing controlled by adjusting p.

Cornered *p*-Norm

#### The norms we consider:

$$\begin{split} \|y\|_{p} &\coloneqq \left(\sum_{i=1}^{k} y_{i}^{p}\right)^{1/p}, \quad p \in [1, \infty) \\ \|y\|_{\infty} &\coloneqq \max_{i=1, \dots, k} y_{i} \\ \|y\|_{p} &\coloneqq \|y\|_{\infty} + \frac{1}{p} \|y\|_{1}, \quad p \in [1, \infty] \end{aligned} \tag{Maximum } (\ell^{\infty} \text{-}) \text{Norm}) \\ \end{split}$$

Weighted version: For any norm and  $\lambda \in \mathbb{Q}^k, \lambda \ge 0, \lambda \ne 0$ :

$$||y||^{\lambda} = ||(\lambda_1 y_1, \lambda_2 y_2, \dots, \lambda_k y_k)||.$$

## **Known Properties**

#### Gearhardt 1979:

- for  $p < \infty$  all compromise solutions are Pareto optimal
- $\bullet\,$  all Pareto optimal solution are a compromise solution, for p big enough



1 Introduction

2 Definitions and Notations

3 Approximation

#### Definition ( $\varepsilon$ -approximate Pareto set)

$$y_i \leq (1+\varepsilon)y'_i \quad \forall i=1,\ldots,k$$



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#### Definition ( $\varepsilon$ -approximate Pareto set)

Let  $\mathcal{Y}_P$  be the Pareto set of a given instance, and let  $\varepsilon > 0$ .  $\mathcal{Y}_{\varepsilon} \subseteq \mathcal{Y}$  is an  $\varepsilon$ -approximate Pareto set if for all  $y \in \mathcal{Y}_P$  there is  $y' \in \mathcal{Y}_{\varepsilon}$  such that

$$y_i \leq (1+\varepsilon)y'_i \quad \forall i=1,\ldots,k$$

#### Theorem (Papadimitriou&Yannakakis,2000)

There always exists an  $\varepsilon$ -approximate Pareto set with size polynomial in |I| and  $1/\varepsilon$ .

Theorem (Papadimitriou&Yannakakis,2000)

There is an efficient algorithm for constructing an  $\varepsilon$ -approximate Pareto set if and only if the GAP problem is tractable.

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### From approximate Pareto sets to approximate CS

Goal: Given an  $\varepsilon$ -approximate Pareto set, find  $(1 + \delta)$ -approximation to the compromise solution

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## An alternative objective

Replace  $\min_{y \in \mathcal{Y}} \|y^* - y\|$  by

 $\max_{y\in\mathcal{Y}}f(y)\;,$ 

where

• level sets are maintained:  $\|y^* - y\| = \|y^* - y'\| \Rightarrow f(y) = f(y')$ 

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$$\Rightarrow f(y) \coloneqq \|y^*\| - \|y^* - y\|$$

## From approximate Pareto sets to approximate CS

#### Theorem

If  $\mathcal{Y}_{\varepsilon}$  is an  $\varepsilon$ -approximate Pareto set for  $\max_{y \in \mathcal{Y}} y$ , then  $\max_{y \in \mathcal{Y}_{\varepsilon}} f(y)$  yields a  $(1 + \delta)$ -approximation to  $\max_{y \in \mathcal{Y}} f(y)$ , for some  $\delta \in \Theta(\varepsilon)$ .

### From approximate CS to approximate Pareto set

Goal: Given FPTAS for CS, solve GAP for given  $y \in \mathbb{R}^k$ ,  $\varepsilon > 0$ .

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## Moving the reference point

#### Solution: Move the reference point



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## Moving the reference point

Solution: Move the reference point ~ super-ideal reference point



 $\dot{y}_1$ 

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### From approximate CS to approximate Pareto set

For 
$$\widetilde{y}^* \ge y^*$$
, set  $\widetilde{f}(y) \coloneqq \|\widetilde{y}^*\| - \|\widetilde{y}^* - y\|$ .

#### Theorem

If there is an FPTAS for  $\max_{y \in \mathcal{Y}} \widetilde{f}(y)$  for any  $\widetilde{y}^* \ge y^*$ , then the GAP problem for  $\max_{y \in \mathcal{Y}} y$  is tractable.

▶ Proof

Approximation

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If  $\mathcal{Y}_{\varepsilon}$  is an  $\varepsilon$ -approximate Pareto set for  $\max_{y \in \mathcal{Y}} y$ , then for any  $\widetilde{y}^* \ge y^*$ ,  $\max_{y \in \mathcal{Y}_{\varepsilon}} \widetilde{f}(y)$  yields a  $(1 + \delta)$ -approximation to  $\max_{y \in \mathcal{Y}} \widetilde{f}(y)$ , for some  $\delta \in \Theta(\varepsilon)$ .

▶ Proof

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# Summary & Outlook

#### approximate Pareto set

ideal point

super-ideal ref pt

#### approximate compromise solutions



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# Summary & Outlook



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# Summary & Outlook



Future work:

- Approximation algorithms
- Heuristics

# Summary & Outlook



#### Future work:

- Approximation algorithms
- Heuristics
- Application

Group of *Multicriteria Analysis and Sustainability*, University of Málaga (Sept–Dec 2012)

# Summary & Outlook



#### Future work:

- Approximation algorithms
- Heuristics
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Thank you for your attention.

Set  $\widetilde{y}^* = C \cdot y$  such that  $\widetilde{y}^* \ge y^*$ . Let y' be a  $(1 + \delta)$ -approximation for  $\max_{y'' \in \mathcal{Y}} \widetilde{f}(y'')$  with  $y' \not\ge y$ .



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$$\forall y'' \in \mathcal{Y} : \widetilde{f}(y'') \leq (1+\delta)\widetilde{f}(y')$$
$$\leq (1+\delta)\widetilde{f}(y)$$
$$= (1+\delta)(\|\widetilde{y}^*\| - \|\widetilde{y}^* - y\|)$$
$$= (1+\delta)(C \cdot \|y\| - (C-1)\|y\|)$$
$$= (1+\delta) \|y\|$$



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$$= (1+\delta)(C \cdot \|y\| - (C-1)\|y\|)$$
  
$$= (1+\delta)\|y\|$$

$$\widetilde{f}((1+\varepsilon)y) = \|\widetilde{y^*}\| - \|\widetilde{y^*} - (1+\varepsilon)y\|$$
$$= C \cdot \|y\| - (C - (1+\varepsilon)) \|y\|$$
$$= (1+\varepsilon) \|y\|$$

Back

Let  $\mathcal{Y}_{\varepsilon}$  be an  $\varepsilon$ -approximate Pareto set,  $y' = \operatorname{argmax}_{y \in \mathcal{Y}_{\varepsilon}} \widetilde{f}(y)$ .  $y^{cs}$  Pareto optimal  $\Rightarrow \exists \overline{y} \in \mathcal{Y}_{\varepsilon}$  with  $\overline{y} \ge (1 - \varepsilon)y^{cs}$ , w.l.o.g.  $\overline{y} = (1 - \varepsilon)y^{cs}$ . We need

$$\|\widetilde{y}^*\| - \|\widetilde{y}^* - \overline{y}\| = \widetilde{f}(\overline{y}) \ge (1 - \delta)\widetilde{f}(y^{\mathsf{cs}}) = (1 - \delta)(\|\widetilde{y}^*\| - \|\widetilde{y}^* - y^{\mathsf{cs}}\|),$$

or equivalently

$$\|\widetilde{y}^* - \overline{y}\| \le \|\widetilde{y}^* - y^{\mathsf{CS}}\| + \delta(\|\widetilde{y}^*\| - \|\widetilde{y}^* - y^{\mathsf{CS}}\|).$$





















