

Fifth Set of Exercises

(due date: May 24, **before** the exercise session)

Notice that this is the last set of exercises in the first half of the semester.

Exercise 23 (4 points). Prove Birkhoff's Theorem by transforming the minimum-weight perfect matching problem in a bipartite graph to a minimum-cost flow problem.

Exercise 24 (4 points). Let $n \in \mathbb{Z}_{>0}$. A matrix $A \in [0, 1]^{n \times n}$ is *doubly stochastic* if the entries in each row and in each column sum up to 1. Show that each doubly stochastic matrix is a convex combination of permutation matrices.

Exercise 25 (4 points). Let $G = (V, E)$ be an undirected graph and $c \in \mathbb{R}^E$. Consider the linear programming relaxation of the minimum-weight perfect matching problem:

$$\begin{aligned} \min \quad & \sum_{e \in E} c_e \cdot x_e \\ \text{s.t.} \quad & x(\delta(v)) = 1 \quad \text{for all } v \in V \\ & x \geq 0 \end{aligned}$$

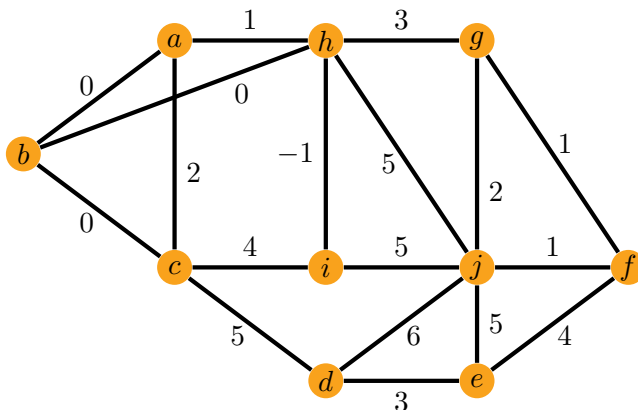
Show that there exists a feasible LP solution if and only if there exists a subset of edges and a set of odd circuits in G such that every node of G is a node of exactly one of the odd circuits or is incident to exactly one of the edges, but not both.

Exercise 26 (4 points). Prove that the optimal value of the following problem is the minimum weight of a perfect matching of G .

$$\begin{aligned} \min \quad & \sum_{e \in E} c_e \cdot x_e \\ \text{s.t.} \quad & x(\delta(v)) = 1 \quad \text{for all } v \in V \\ & x(\gamma(S)) \leq (|S| - 1)/2 \quad \text{for all } S \subseteq V, |S| \geq 3, |S| \text{ odd} \\ & x \geq 0 \end{aligned}$$

Exercise 27 (4 points). Let M be a perfect matching of G and define the "cost" of an M -alternating circuit C of G to be $c(E(C) \setminus M) - c(E(C) \cap M)$. Prove that M is of minimum weight with respect to c if and only if there is no M -alternating circuit of negative cost.

Exercise 28 (tutorial session). Find a minimum-weight perfect matching and an optimal solution of the dual problem for the following weighted graph:



Exercise 29 (tutorial session). Let $G = (V, E)$ be an undirected graph, let $c \in \mathbb{R}^E$, and let k be a positive integer. Show how the problem of finding a minimum-weight matching having cardinality k can be reduced to a minimum-weight perfect matching problem.

Exercise 30 (tutorial session). Show that there exists a non-bipartite graph with the property that, for every $c \in \mathbb{R}^E$, the optimal value of the linear program in Exercise 25 is equal to the minimum weight of a perfect matching of G .

Exercise 31 (tutorial session). Let $G = (V, E)$ be an undirected graph. Prove that G has a perfect matching if and only if there is an $x \in \mathbb{R}^E$ with:

$$\begin{aligned}
 x(\delta(v)) &= 1 && \text{for all } v \in V \\
 x(\delta(S)) &\geq 1 && \text{for all } S \subseteq V, |S| \text{ odd} \\
 x &\geq 0
 \end{aligned}$$