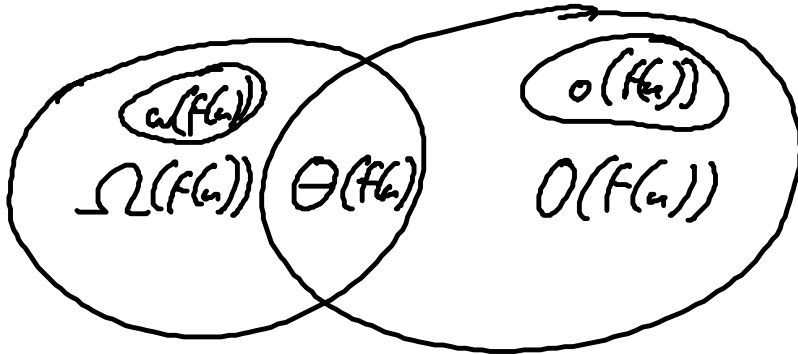


CoMa-Übung 12, 04.02.13

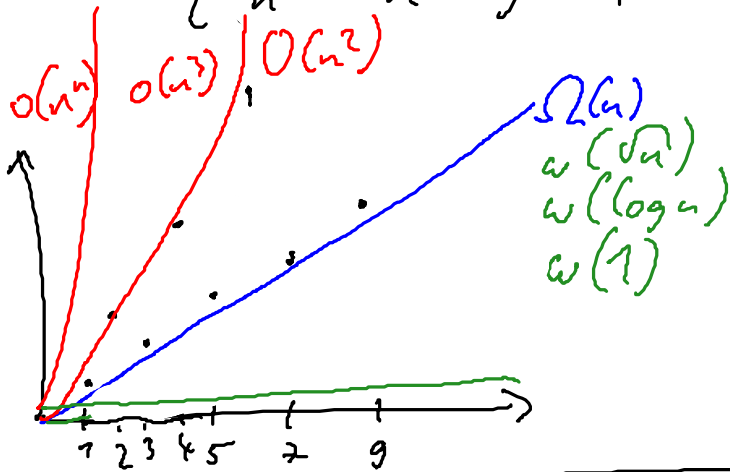
- Anmeldung zu Proberücksprachen
- " = Rücksprachen
- Wiederholung



$$f(n) = \begin{cases} n^2 & n \text{ gerade} \\ n & n \text{ ungerade} \end{cases}$$

$$f(n) \in O(n^2) \checkmark$$

$$f(n) \in \Omega(n) \checkmark$$



Asymptotisches Wachstum

Asymptotisch langsamer

Konstante Funktionen  $f(n) = 15 \rightarrow \Theta(1)$   $\rightarrow$  z.B. Zugriff auf Array-Element

Logarithmische Funktionen  $f(n) = \log n \rightarrow \Theta(\log n)$   $\rightarrow$  z.B. Binäre Suche

Wurzel(-Funktion)  $f(n) = \sqrt{n} = n^{\frac{1}{2}} \rightarrow \Theta(n^{\frac{1}{2}})$

Polynome  $f(n) = \sum_{i=1}^k a_i \cdot n^i \rightarrow \Theta(n^k)$   $n^3 + n^2 + n \rightarrow \Theta(n^3)$

$$= \sum a_i \cdot n^i \quad I \in \mathbb{R}$$

$$n^2 + \sqrt{n} \rightarrow \Theta(n^2)$$

$$i \in I \rightarrow \Theta(n^{\max i})$$

Exponentielle Funktionen  $f(n) = 2^n \rightarrow$  Vollständige Enumerierung  
 $\rightarrow \Theta(2^n)$

Kombinatorische Funktionen  $f(n) = n! \rightarrow$  Alle Permutationen  
enumerieren

Asymptotisch  
schneller

$$f \in \Theta(n^2) \Leftrightarrow f \in \Omega(n^2) \wedge f \in O(n^2)$$

von ☺

0, public boolean recurse (Strings)

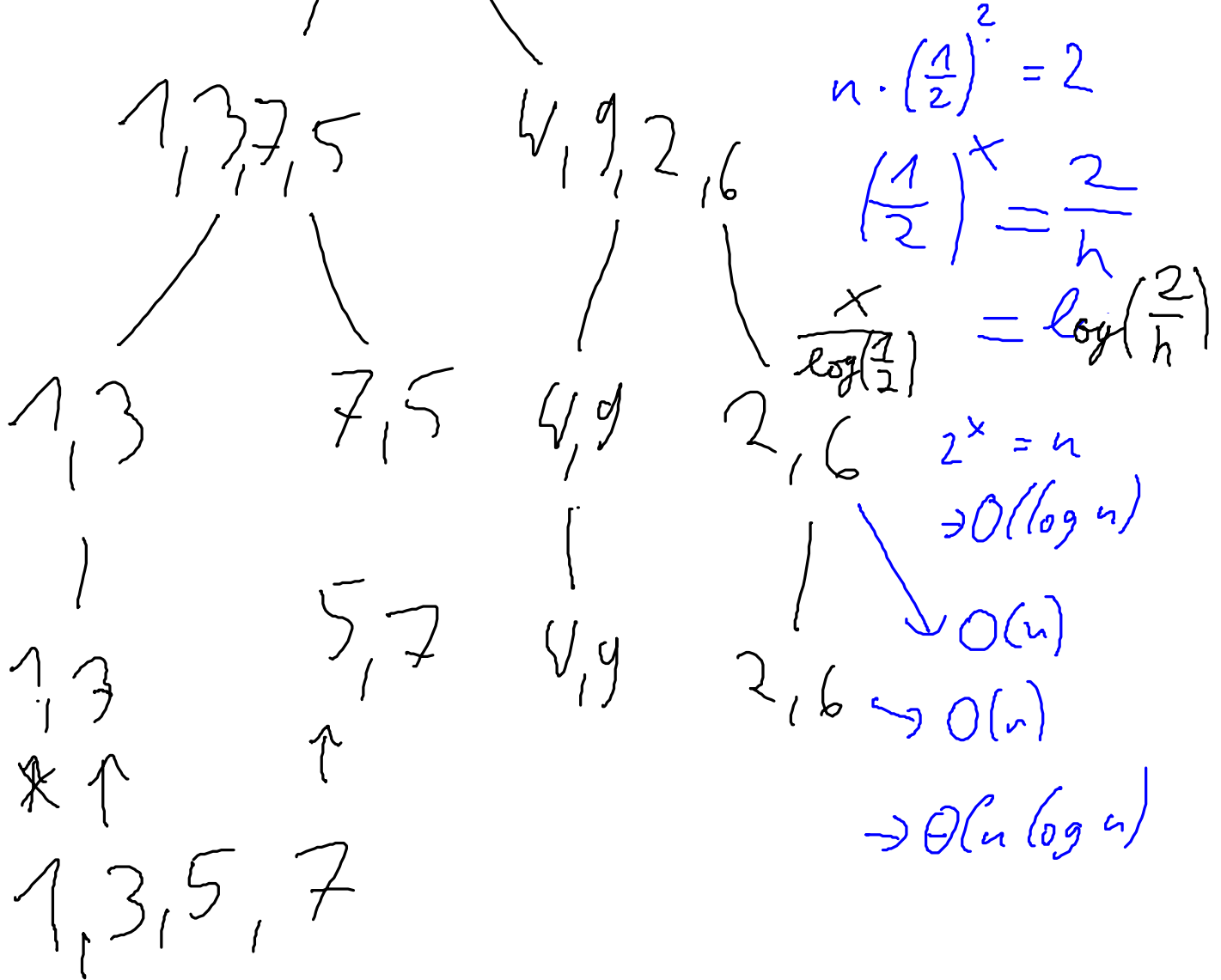
```
{ if (s.length() == 0) return true;
```

else

```
{ if (s.charAt(0) == s.charAt(s.length() - 1))  
return (recurse (s.substring(1, s.length() - 2));
```

```
else return false; } }
```

1, 3, 7, 5, 4, 9, 2, 6  
| | | | | |  
| | | | | x



```

public class arr {
    private double[][] arr;
    public arr(double[][] i)
    {
        *
        for (int a = 0; a < i.length; a++)
    }
}

```

```

    } for (int b = 0; b < i[0].length; b++)
        { varis[a][b] = i[a][b]; } }

```

\* = varis = new double [i.length][i[0].length];

```

public ans ()

```

```

{ return

```

```

    (this.ans (new double [2][2]); } }

```

```

public class matr extends ans implements matr

```

```

{ public matr (double [][] d) *

```

```

    { for (int f = 0; f < d.length; f++)

```

```

        { for (int g = 0; g < d[0].length; g++)

```

```
{ if (d[f][g] != d[g][f])  
    throw new Exception(); }
```

\* Throws Exception

```
super(d); }
```

```
public interface math
```

```
{ }
```

$f(1) = 3; f(2) = 2; f(3) = 1;$   
 $f(n) = 3 \cdot f(n-1) + f(n-2) \cdot f(n-3)$

```
public int recur(int n)
```

```
{ if (n == 1) return 3;
```

```

else if (n == 2) return 2;
else if (n == 3) return 1;
else return (return(n-1) + return(n-2) * return(n-3));

```

$\} \quad 2^n \in O(n^2)$

$$f(n) \in O(g(n)) \Leftrightarrow \exists c \in \mathbb{R} \exists n_0, (c \cdot f(n) \leq g(n))$$

$$\rightarrow c \cdot 2^n \leq n^2 \quad \forall c \in \mathbb{R} \\ \forall n \geq n_0$$

$\forall n \geq n_0$   
 $\rightarrow$  ab dem  $n_0$  dominiert das asymptotische Wachstum

TCJ

ArrayList < T >  $\frac{n^3 + 1}{n^2} = n + \frac{1}{n^2}$

(class < T >

$$n + \frac{1}{n^2} \in O(n)$$

String.class

Integer.class  
 LinkedList.class

T.class

public < T > TCJ createArray (Class < T > c, int (length) l